# Analyzing an Agent-Based Model for House-Hunting in Ant Colonies

Emily Zhang, Jiajia Zhao, Nancy Lynch Massachusetts Institute of Technology

# The House-Hunting Algorithm

- Model captures biological measurements from empirical studies (Zhao, Lynch, Pratt, 2020)
  - Bio-plausible
  - Predicts less studied behaviors of ants
- **Challenge:** Lack theoretical bounds on the running time of the algorithm



Phase transitions of an ant

# Simplifications



setf.nome\_nest = \_nome\_nest # default to 0
self.candidate\_nest = \_candidate\_nest
self.transitioned = \_transitioned
self.location = \_home\_nest
self.phase = "E" # how much ant is committed to the candidate nest
self.old\_candidate\_nest = -1 # only used for reject action
self.terminate = 0

# After



# Adjustable Parameters

Parameter	Value	Source
quality coefficient $\mu_q$	0.25	trial-and-error from $[13]$
population coefficient $\mu_p$	0.35	trial-and-error from $[13]$
quorum threshold $\theta$	0.15	[9, 2]
search constant $c_s$	0.025	trial-and-error from [13]
follow constant $c_f$	0.4	[4, 8]
lead forward constant $c_\ell$	0.6	trial-and-error from $[13]$
transport constant $c_t$	0.7	[10]
$\lambda$	8	trial-and-error from $[13]$

# **State Transitions**



$$\begin{aligned} &\Pr[\mathbf{u} = \operatorname{advance} \mid a.state = \operatorname{At} \operatorname{Nest}_i] = \left(1 + e^{-\lambda \left(\mu_q \cdot q + \mu_p \cdot \frac{p}{n}\right)}\right)^{-1} \text{ for } i \in \{\operatorname{E}, \operatorname{C}, \operatorname{T}\} \\ &\Pr[\mathbf{u} = \operatorname{advance} \mid a.state = \operatorname{Search}_i] = c_s \cdot \left(1 + e^{-\lambda \left(\mu_q \cdot (q' - q) + \mu_p \cdot \frac{p' - p}{n}\right)}\right)^{-1} \text{ for } i \in \{\operatorname{E}, \operatorname{C}, \operatorname{T}\} \\ &\Pr[\mathbf{u} = \operatorname{advance} \mid a.state = \operatorname{Quorum} \operatorname{Sensing}] = \begin{cases} 1 & \text{if quorum has been met} - \text{ that is, } p_a > \theta \cdot n_a \\ & \text{and } a.location \text{ has not dropped out of competition} \\ 0 & \text{otherwise} \end{cases} \\ &\Pr[\mathbf{u} = \operatorname{advance} \mid a.state = \operatorname{Transport}] = c_t \end{aligned}$$

q'

$$\Pr[\mathbf{u} = \text{advance} \mid a.state = \text{Lead Forward}] = \begin{cases} c_{\ell} & \text{if } q > q' \\ 0 & \text{otherwise} \end{cases}$$

# House-Hunting Algorithm

Algorithm 1: One Round of the HOUSEHUNTING Algorithm		
<b>1</b> $M$ : a set of ants, initially $\emptyset$		
2 for $i = 1$ to $n_a$ do		
3 if $a_{P(i)} \notin M$ then		
4	$action, n' := $ <b>select_action</b> $(a_{P(i)})$	
5	$a' := \mathbf{select\_ant}(\mathbf{a_{P(i)}}, \mathbf{n'}, \mathbf{action})$	
6	if $a' \in M$ then	
7	$  a' \leftarrow null$	
8	$transition(a_{P(i)}, a', n', action)$	
9	$M := M \cup \left\{ a_{P(i)} \right\} \cup \left\{ a' \right\}$	

# A Lower Bound on Number of Rounds Required

Proof inspired by Ghaffari, Musco, Radeva, Lynch, 2015.

#### Method:

- Lower bound the probability that a constant fraction of the ants goes to the new nest on any given round.
- Chernoff Bound

#### **Result**:

• An algorithm with *n* ants requires  $\Omega(\log n)$  rounds to converge with high probability.

**Theorem 4.4.** If the quorum threshold satisfies  $1 - \frac{a(\epsilon)}{n_a} < \theta < \frac{a(\epsilon)}{n_a}$ , then  $\mathbb{E}[R_{\epsilon}] = O(\log n)$ .

$$n_a =$$
 the number of active ants  
 $a(\epsilon) \approx n_a \left(\frac{1-\epsilon}{1+e^{-\lambda(\mu_q(q_1-q_0)-\mu_p)}}\right)$ 

**Theorem 4.4.** If the quorum threshold satisfies  $1 - \frac{a(\epsilon)}{n_a} < \theta < \frac{a(\epsilon)}{n_a}$  then  $\mathbb{E}[R_{\epsilon}] = O(\log n)$ .

 $n_a =$  the number of active ants  $a(\epsilon) \approx n_a \left(\frac{1-\epsilon}{1+e^{-\lambda(\mu_q(q_1-q_0)-\mu_p)}}\right)$ 

**Lower bound**: transports to inferior nest taper off **Upper bound**: transports to superior nest begin in O(log n) rounds

**Theorem 4.4.** If the quorum threshold satisfies 
$$1 - \frac{a(\epsilon)}{n_a} < \theta < \frac{a(\epsilon)}{n_a}$$
, then  $\mathbb{E}[R_{\epsilon}] = O(\log n)$ .

$$n_a =$$
 the number of active ants  
 $a(\epsilon) \approx n_a \left(\frac{1-\epsilon}{1+e^{-\lambda(\mu_q(q_1-q_0)-\mu_p)}}\right)$ 

**Lower bound**: transports to inferior nest taper off **Upper bound**: transports to superior nest begin in O(log n) rounds

$$\begin{split} \lambda &= 8, \ \mu_q = .25, \ \mu_p = .35, \ q_1 = 3, \ q_0 = 0, \ \epsilon = .00001 \\ \theta &\in (.0392, .9608) \quad \text{(reasonable bounds)} \\ \theta &\approx .15 \end{split}$$

**Theorem 4.4.** If the quorum threshold satisfies 
$$1 - \frac{a(\epsilon)}{n_a} < \theta < \frac{a(\epsilon)}{n_a}$$
 then  $\mathbb{E}[R_{\epsilon}] = O(\log n)$ .  
 $n_a = \text{the number of active ants}$   
 $a(\epsilon) \approx n_a \left(\frac{1-\epsilon}{1+e^{-\lambda(\mu_q(q_1-q_0)-\mu_p)}}\right)$   
Lower bound: transports to inferior nest taper off  
Upper bound: transports to superior nest begin in O(log n)  
rounds

#### **Future Work**:

1. Generalize this result to environments with multiple nests

$$1 - \frac{a(\epsilon)}{n_a} < \frac{a(\epsilon)}{n_a} \implies q_1 - q_0 > \frac{\mu_p}{\mu_q}$$

Test this prediction about the relative qualities of nests in real experiments with ants

# Implications

- Gain a better understanding of the biological behavior of ants.
- Studying biologically-inspired algorithms can help engineer better distributed computer systems.
  - Robot swarms

# Thank you. Questions?